

# Mathematical Logics

## FOL: Reasoning as deduction

Fausto Giunchiglia and Mattia Fumagalli

University of Trento



*\*Originally by Luciano Serafini and Chiara Ghidini  
Modified by Fausto Giunchiglia and Mattia Fumagalli*

1. Reasoning problems (recap)
2. Hilbert systems (VAL – forward chaining)
3. Tableaux systems ( (un)-SAT – backward chaining)
4. Correctness and completeness of Tableau
5. Examples
6. Termination
7. Countermodels

## Exercise

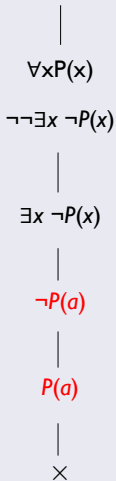
Show with the method of semantic tableaux that the following formulas are valid:

- $\forall x P(x) \supset \neg \exists x \neg P(x)$
- $\forall x (P(x) \vee A) \supset (\forall x P(x) \vee A)$  when  $x$  is not free in  $A$
- $\exists x (P(x) \supset \forall x P(x))$
- $\exists x \forall y P(x, y) \supset \forall y \exists x P(x, y)$

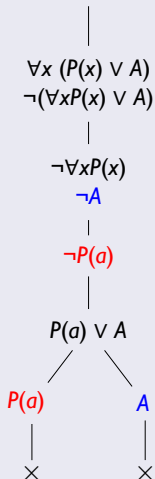
# Practicing with Tableaux

## Solution

$$\neg(\forall x P(x) \supset \neg \exists x \neg P(x))$$



$$\neg(\forall x (P(x) \vee A) \supset (\forall x P(x) \vee A))$$



# Practicing with Tableaux

## Solution

$$\neg \exists x (P(x) \supset \forall x P(x))$$

$$\neg (P(a) \supset \forall x P(x))$$

$$P(a)$$

$$\neg \forall x P(x)$$

$$\neg P(b)$$

$$\neg (P(b) \supset \forall x P(x))$$

$$P(b)$$

$$\neg \forall x P(x)$$

×

$$\neg (\exists x \forall y P(x, y) \supset \forall y \exists x P(x, y))$$

$$\exists x \forall y P(x, y)$$

$$\neg \forall y \exists x P(x, y)$$

$$\forall y P(a, y)$$

$$\neg \exists x P(x, b)$$

$$P(a, b)$$

$$\neg P(a, b)$$

×

## Exercize

Give tableau proofs for the following logical consequences:

- $\forall x.P(x) \vee \forall x.Q(x) \models \neg \exists x (\neg P(x) \wedge \neg Q(x))$
- $\models \exists x.(P(x) \vee Q(x)) \equiv \exists x.P(x) \vee \exists x.Q(x)$

# Prove the following FOL properties of quantifiers

## Proposition

The following formulas are valid

- $\forall x (\varphi(x) \wedge \psi(x)) \equiv \forall x \varphi(x) \wedge \forall x \psi(x)$
- $\exists x (\varphi(x) \vee \psi(x)) \equiv \exists x \varphi(x) \vee \exists x \psi(x)$
- $\forall x \varphi(x) \equiv \neg \exists x \neg \varphi(x)$
- $\forall x \exists x \varphi(x) \equiv \exists x \varphi(x)$
- $\exists x \forall x \varphi(x) \equiv \forall x \varphi(x)$

## Proposition

The following formulas are not valid (prove the correct direction)

- $\forall x (\varphi(x) \vee \psi(x)) \equiv \forall x \varphi(x) \vee \forall x \psi(x)$
- $\exists x (\varphi(x) \wedge \psi(x)) \equiv \exists x \varphi(x) \wedge \exists x \psi(x)$
- $\forall x \varphi(x) \equiv \exists x \varphi(x)$
- $\forall x \exists y \varphi(x, y) \equiv \exists y \forall x \varphi(x, y)$

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