

Mathematical Logics

Description Logic: Tbox and Abox

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Modified by Fausto Giunchiglia and Mattia Fumagalli*

1. Families of Description Logics
2. TBOX: syntax and semantics
3. TBOX: terminology
4. TBOX: reasoning
5. ABOX: syntax and semantics
6. ABOX: reasoning
7. Closed World Assumption (CWA) and Open World Assumption (OWA)

Given two class-propositions P and Q , we have the following reasoning problems

□ Satisfiability w.r.t. T $T \models P$?

□ Subsumption $T \models P \sqsubseteq Q$? $T \models Q \sqsubseteq P$?

□ Equivalence $T \models P \sqsubseteq Q$ and $T \models Q \sqsubseteq P$?

□ Disjointness $T \models P \sqcap Q \sqsubseteq \perp$?

Satisfiability with respect to a TBox T

A concept P is **satisfiable w.r.t. a terminology T** , if there exists (or for all) an interpretation I with $I \models \theta$ for all $\theta \in T$, and such that $I \models P$, namely $I(P)$ is not empty

In this case we also say that **I is a model** for P

In other words, the interpretation I not only satisfies P , but also complies with all the constraints in T

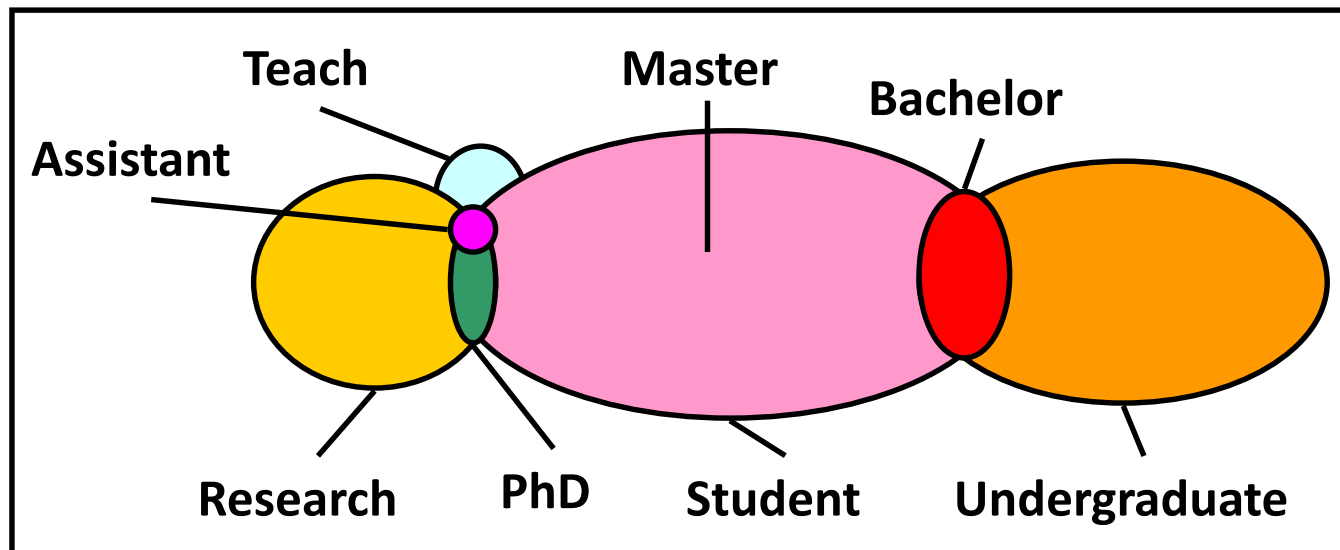
Satisfiability with respect to a TBox T

Suppose we describe the students in a course:

Undergraduate	$\sqsubseteq \neg \text{Teach}$
Bachelor	$\equiv \text{Student} \sqcap \text{Undergraduate}$
Master	$\equiv \text{Student} \sqcap \neg \text{Undergraduate}$
PhD	$\equiv \text{Master} \sqcap \text{Research}$
Assistant	$\equiv \text{PhD} \sqcap \text{Teach}$

TBox T

The TBox is satisfiable. A possible model is:



In this model the two concepts **Bachelor** and **Assistant** are satisfiable w.r.t. T, while the concept **Assistant \sqcap Bachelor** is not.

TBox reasoning: subsumption

Let T be a TBox. A **Subsumption** problem (with respect to T): is defined as follows

$$T \models P \sqsubseteq Q \quad (P \sqsubseteq_T Q)$$

A concept P is subsumed by a concept Q with respect to T if $I(P) \subseteq I(Q)$ for every model I of T

NOTE: **Subsumption** reduces to entailment and validity (when T empty)

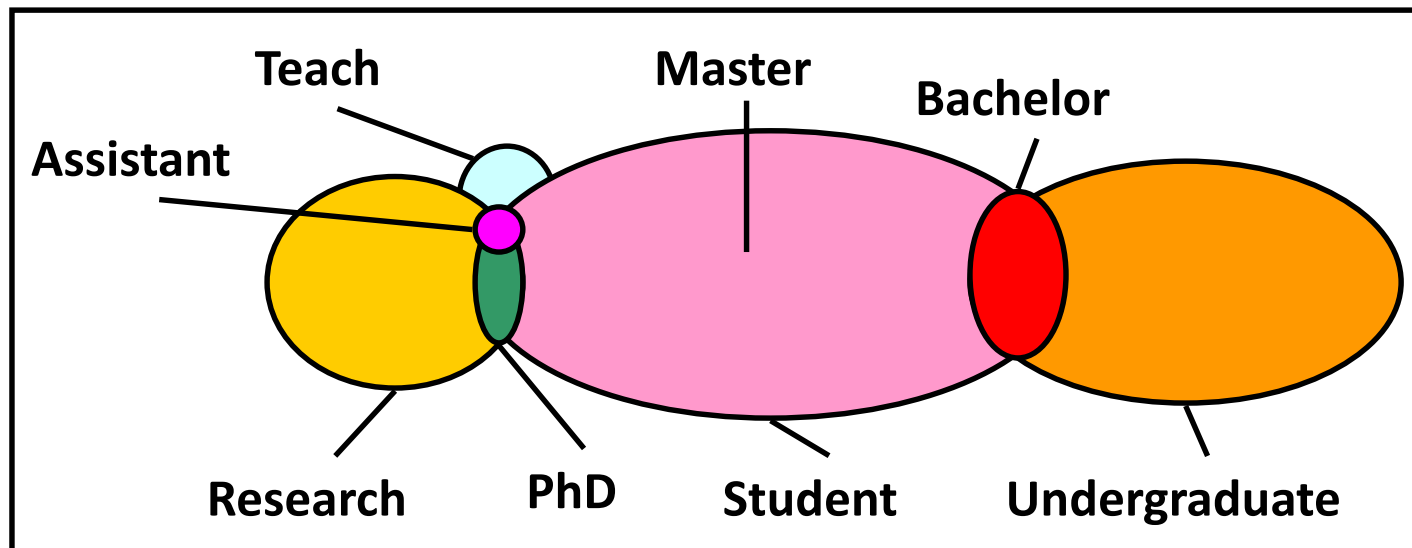
Subsumption with respect to a TBox T (I)

Suppose we describe the students in a course:

Undergraduate	$\sqsubseteq \neg \text{Teach}$
Bachelor	$\equiv \text{Student} \sqcap \text{Undergraduate}$
Master	$\equiv \text{Student} \sqcap \neg \text{Undergraduate}$
PhD	$\equiv \text{Master} \sqcap \text{Research}$
Assistant	$\equiv \text{PhD} \sqcap \text{Teach}$

TBox T

T \models PhD \sqsubseteq Student



Subsumption with respect to a TBox T (2)

PhD \sqsubseteq Student

Proof:

PhD

\equiv Master \sqcap Research

\equiv (Student $\sqcap \neg$ Undergraduate) \sqcap Research

\sqsubseteq Student

TBox reasoning: equivalence

Let T be a TBox. An **Equivalence** problem (with respect to T) is defined as follows:

$$(T \models P \equiv Q) \quad (P \equiv_T Q)$$

Two concepts P and Q are equivalent with respect to T if $I(P) = I(Q)$ for every model I of T .

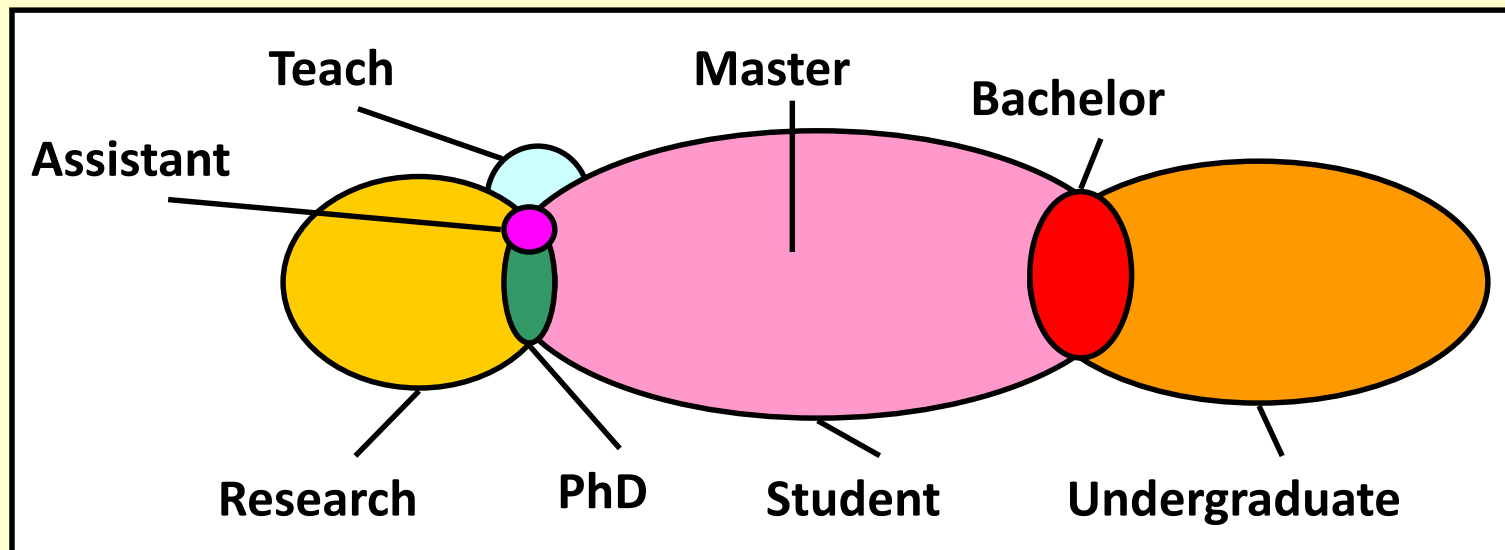
Equivalence with respect to a TBox \mathcal{T} (I)

Suppose we describe the students in a course:

Undergraduate	$\sqsubseteq \neg \text{Teach}$
Bachelor	$\equiv \text{Student} \sqcap \text{Undergraduate}$
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Assistant	$\equiv \text{PhD} \sqcap \text{Teach}$

TBox \mathcal{T}

$\mathcal{T} \models \text{Student} \equiv \text{Bachelor} \sqcup \text{Master}$



Equivalence with respect to a TBox T (2)

Student \equiv Bachelor \sqcup Master

Proof:

Bachelor \sqcup Master

\equiv (Student \sqcap Undergraduate) \sqcup Master

\equiv (Student \sqcap Undergraduate) \sqcup (Student \sqcap \neg Undergraduate)

\equiv Student \sqcap (Undergraduate \sqcup \neg Undergraduate)

\equiv Student \sqcap T

\equiv Student

Let T be a TBox. A **Disjointness** problem (with respect to T) is defined as follows:

$$T \models P \sqcap Q \sqsubseteq \perp \quad (P \sqcap Q \sqsubseteq_T \perp)$$

Two concepts P and Q are disjoint with respect to T if their intersection is empty, $I(P) \cap I(Q) = \emptyset$, for every model I of T .

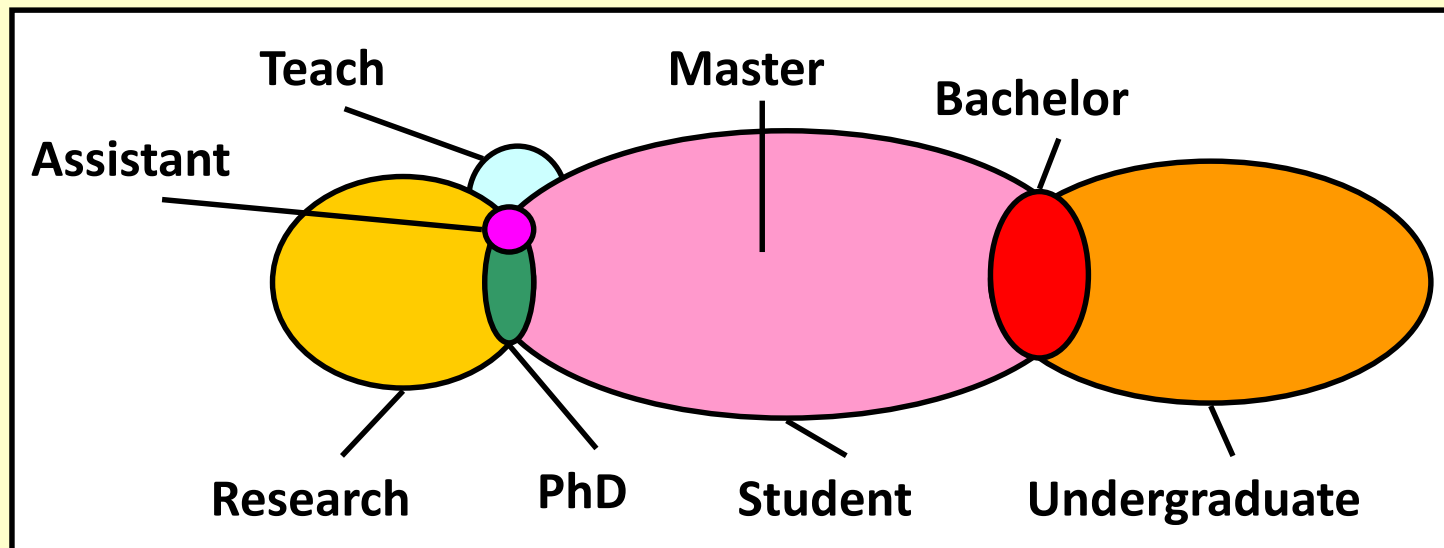
Disjointness with respect to a TBox \mathcal{T} (I)

Suppose we describe the students in a course:

Undergraduate	$\sqsubseteq \neg \text{Teach}$
Bachelor	$\equiv \text{Student} \sqcap \text{Undergraduate}$
Master	$\equiv \text{Student} \sqcap \neg \text{Undergraduate}$
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Assistant	$\equiv \text{PhD} \sqcap \text{Teach}$

TBox \mathcal{T}

$\mathcal{T} \models \text{Undergraduate} \sqcap \text{Assistant} \sqsubseteq \perp$



Disjointness with respect to a TBox \mathcal{T} (2)

It can be proved showing that:

$$\mathcal{T} \models \text{Undergraduate} \sqcap \text{Assistant} \sqsubseteq \perp$$

Proof:

Undergraduate \sqcap Assistant

$$\sqsubseteq \neg \text{Teach} \sqcap \text{Assistant}$$

$$\equiv \neg \text{Teach} \sqcap \text{PhD} \sqcap \text{Teach}$$

$$\equiv \perp \sqcap \text{PhD}$$

$$\equiv \perp$$

Exercise

Suppose we describe the students in a course:

Undergraduate	$\sqsubseteq \neg \text{Teach}$
Bachelor	$\equiv \text{Student} \sqcap \text{Undergraduate}$
Master	$\equiv \text{Student} \sqcap \neg \text{Undergraduate}$
PhD	$\equiv \text{Master} \sqcap \text{Research}$
Assistant	$\equiv \text{PhD} \sqcap \text{Teach}$

TBox T

Is $\text{Bachelor} \sqcap \text{PhD}$ satisfiable?

NO!

Consider the following propositions:

Assistant, Student, Bachelor, Teach, PhD, Master \sqcap Teach

- Which pairs are subsumed/supersumed?
- Which pairs are disjoint?

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