

Mathematical Logics

Description Logic: Tableaux

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1. Idea: DL is a MultiModal Modal Logic
2. DL reasoning as MultiModal SAT reasoning
3. Examples: TBOX reasoning
4. Examples: ABOX reasoning – DL as a query language

DL as Query Language - intuition

TBox = {}

ABox = {Child(John, Mary), Female(Mary)}

NL Query: **Who are the individuals having only female children?**

DL Query: $T, A \models \forall \text{Child.Female}$

Answer: {John}

□ ABox services are generally applied to resolve a query

How to use ABox Reasoning Services

ABox Service	Description	Query
Instance retrieval	Given a concept C , retrieve all the instances a which satisfy C w.r.t. the ABox A .	$A \models C$
Instance checking	Check whether an assertion $C(a)$ is entailed by the ABox, i.e. check whether a belongs to C .	$A \models C(a)$ $A \models R(a,b)$

NOTE: this means that before answering we need to expand the ABox (w.r.t. the TBox) and reason on the identified model (see before)

Answering Queries via instance checking (I)

TBox = {Horse \sqsubseteq Animal, Mule \sqsubseteq Animal}

ABox = {Horse(Furia), Parent(Speedy, Furia)}

NL Query: Is Furia an animal?

DL Query: T, A \models Animal(Furia)

YES, in fact the ABox can be expanded as follows:

ABox = {Horse(Furia), Animal(Furia), Parent(Speedy, Furia)}

Answering Queries via instance checking (II)

TBox = {Horse \sqsubseteq Animal \sqcap \neg Mule, Mule \sqsubseteq Animal}

ABox = {Horse(Furia), Parent(Speedy, Furia)}

NL Query: Is Furia a mule?

DL Query: T, A \models Mule(Furia)

NO, in fact the ABox can be expanded as follows:

ABox = {Horse(Furia), Animal(Furia), \neg Mule(Furia),
Parent(Speedy, Furia)}

Answering Queries via instance checking (II)

TBox = {Horse \sqsubseteq Animal, Mule \sqsubseteq Animal}

ABox = {Horse(Furia), Parent(Speedy, Furia)}

NL Query: **Is Furia a mule?**

DL Query: $T, A \models \text{Mule}(\text{Furia})$

NO (BY CLOSED WORLD ASSUMPTION), in fact the ABox can be expanded as follows:

ABox = {Horse(Furia), Animal(Furia), Parent(Speedy, Furia)}

If we drop closed world assumption the answer should be:

I DO NOT KNOW

Answering Queries via instance retrieval: Tableaux (I)

TBox = {Horse \sqsubseteq Animal, Mule \sqsubseteq Animal}

ABox = {Horse(Speedy), Horse(Furia), Parent(Speedy, Furia)}

NL Query: Is there any animal which is not a horse nor a mule, and is parent of a horse?

DL Query: $T, A \models \exists \text{Parent.Horse} \sqcap \neg (\text{Horse} \sqcap \text{Mule})$
i.e. is the formula satisfiable?

Answering Queries via instance retrieval: Tableaux (I)

TBox = {Horse \sqsubseteq Animal, Mule \sqsubseteq Animal}

ABox = {Horse(Speedy), Horse(Furia), Parent(Speedy, Furia)}

Is $\exists \text{Parent.Horse} \sqcap \neg (\text{Horse} \sqcap \text{Mule})$ satisfiable?

\sqcap -rule $A' = \{ \exists \text{Parent.Horse}(x), \neg(\text{Horse} \sqcap \text{Mule})(x) \}$

\exists -rule $A' = \{ \text{Horse}(Furia), \text{Parent}(\text{Speedy}, \text{Furia}), (\neg \text{Horse} \sqcup \neg \text{Mule})(x) \}$

\sqcup -rule $A' = \{ \text{Horse}(Furia), \text{Parent}(\text{Speedy}, \text{Furia}), \neg \text{Horse}(Furia) \}$
inconsistent!

or

$A' = \{ \text{Horse}(Furia), \text{Parent}(\text{Speedy}, \text{Furia}), \neg \text{Mule}(Furia) \}$
consistent!

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