

# Computational Logic

## L7.X.1 PL Exercises

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# Outline

- Informal 2 Formal
  - Rules of the thumb
  - Examples
- Reasoning
  - Truth Tables
  - Tableaux

# Formalizing NL

- It is the case that P:  
 $P$
- It is not the case that P:  
 $\neg P$
- P and Q. P but Q. Although P, Q:  
 $(P \wedge Q)$
- P or Q. Either P or Q:  
 $(P \vee Q)$
- P if and only if Q:  
 $(P \leftrightarrow Q)$
- If P, then Q:  
 $(P \rightarrow Q)$
- P if Q:  
 $(Q \rightarrow P)$
- Neither P nor Q:  
 $\neg (P \vee Q)$   
 $\neg P \wedge \neg Q$
- It is not the case that both P and Q:  
 $\neg (P \wedge Q)$
- Both not P and not Q:  
 $(\neg P \wedge \neg Q)$

# Formalizing NL

- It is the case that P:  
P
- It is not the case that P:  
not P
- P and Q. P but Q. Although P, Q:  
(P and Q)
- P or Q. Either P or Q:  
(P or Q)
- P if and only if Q:  
(P iff Q)
- If P, then Q:  
(P imp Q)
- P if Q:  
(Q imp P)
- Neither P nor Q:  
not (P or Q)  
not P and not Q
- It is not the case that both P and Q:  
not (P and Q)
- Both not P and not Q:  
(not P and not Q)

# Meaning of Implication

“The meaning of the implication operator ( $\rightarrow$ ) may appear unintuitive, since we must get used to the notion that “falsehood implies everything.” We should not confuse ( $\rightarrow$ ) with causation. That is,  $p \rightarrow q$  may be true, yet  $p$  does not “cause”  $q$  in any sense. For example, let  $p$  be “it is raining,” and  $q$  be “Sue takes her umbrella.” We might assert that  $p \rightarrow q$  is true. It might even appear that the rain is what caused Sue to take her umbrella. However, it could also be true that Sue is the sort of person who doesn’t believe weather forecasts and prefers to carry an umbrella at all times”

(<http://infolab.stanford.edu/~ullman/focs/ch12.pdf>)

# Meaning of Implication

- Use implication to restrict conditions: “In the situations in which ... then ...”
- Probably simpler to grasp if you reason with the negation (It is not the case that the consequence is false and the premiss is true)

# Formalizing NL

Let us consider a propositional language where **ph** means "Paola is happy", **pp** means "Paola paints a picture", and **rh** means "Renzo is happy".

Formalize the following sentences:

- "if Paola is happy and paints a picture then Renzo is not happy"  
 $(ph \wedge pp) \rightarrow \neg rh$
- "if Paola is happy, then she paints a picture"  
 $ph \rightarrow pp$
- "Paola is happy only if she paints a picture"  
 $\neg (ph \wedge \neg pp)$  or, equivalently,  $ph \rightarrow pp$

# Formalizing NL

- "If Davide comes to the party then Bruno and Carlo come too"
- "Carlo comes to the party only if Angelo and Bruno do not come"
- "If Davide comes to the party, then, if Carlo doesn't come then Angelo comes"
- "Carlo comes to the party provided that Davide doesn't come, but, if Davide comes, then Bruno doesn't come"
- "A necessary condition for Angelo coming to the party, is that, if Bruno and Carlo aren't coming, Davide comes"
- "Angelo, Bruno and Carlo come to the party if and only if Davide doesn't come, but, if neither Angelo nor Bruno come, then Davide comes only if Carlo comes"

# Formalizing NL

- "If Davide comes to the party then Bruno and Carlo come too"  
 $D \rightarrow (B \wedge C)$
- "Carlo comes to the party only if Angelo and Bruno do not come"  
 $C \rightarrow (\neg A \wedge \neg B)$
- "If Davide comes to the party, then, if Carlo doesn't come then Angelo comes"  
 $D \rightarrow (\neg C \rightarrow A)$
- "Carlo comes to the party provided that Davide doesn't come, but, if Davide comes, then Bruno doesn't come"  
 $(C \rightarrow \neg D) \wedge (D \rightarrow \neg B)$
- "A necessary condition for Angelo coming to the party, is that, if Bruno and Carlo aren't coming, Davide comes"  
 $A \rightarrow (\neg B \wedge \neg C \rightarrow D)$
- "Angelo, Bruno and Carlo come to the party if and only if Davide doesn't come, but, if neither Angelo nor Bruno come, then Davide comes only if Carlo comes"  
 $(A \wedge B \wedge C \leftrightarrow \neg D) \wedge (\neg A \wedge \neg B \rightarrow (D \rightarrow C))$

# List the models for the propositions

□ List the models for the following formulas:

1.  $A \wedge \neg B$

2.  $(A \wedge B) \vee (B \wedge C)$

3.  $(A \vee B) \rightarrow C$

4.  $(\neg A \leftrightarrow B) \leftrightarrow C$

**MODEL**



A	B	$A \wedge \neg B$
T	T	F
T	F	T
F	T	F
F	F	F

# Truth Tables and Validity (1)

Use the truth tables method to determine whether  $(p \rightarrow q) \vee (p \rightarrow \neg q)$  is valid.

$p$	$q$	$p \rightarrow q$	$\neg q$	$p \rightarrow \neg q$	$(p \rightarrow q) \vee (p \rightarrow \neg q)$
T	T	T	F	F	T
T	F	F	T	T	T
F	T	T	F	T	T
F	F	T	T	T	T

The formula is valid since it is satisfied by every interpretation.

# Truth Tables and Validity (2)

Use the truth tables method to determine whether  $(\neg p \vee q) \wedge (q \rightarrow \neg r \wedge \neg p) \wedge (p \vee r)$  (denoted with  $F$ ) is valid.

$p$	$q$	$r$	$\neg p \vee q$	$\neg r \wedge \neg p$	$q \rightarrow \neg r \wedge \neg p$	$(p \vee r)$	$F$
T	T	T	T	F	F	T	F
T	T	F	T	F	F	T	F
T	F	T	F	F	T	T	F
T	F	F	F	F	T	T	F
F	T	T	T	F	F	T	F
F	T	F	T	T	T	F	F
F	F	T	T	F	T	T	T
F	F	F	T	T	T	F	F

There exists an interpretation satisfying  $F$ , thus  $F$  is satisfiable.

# Problem formalization (2)

Brown, Jones, and Smith are suspected of a crime. They testify as follows:

- Brown: “Jones is guilty and Smith is innocent”.
- Jones: “If Brown is guilty then so is Smith”.
- Smith: “I'm innocent, but at least one of the others is guilty”.

Formalization:

1. Choose an appropriate propositional representation for the sentences: Let  $B$ ,  $J$ , and  $S$  be the statements “Brown is guilty”, “Jones is guilty”, and “Smith is guilty”, respectively.
2. Express the sentences (depositions) of each suspect as a PL formula. The three statements can be expressed as  $J \wedge \neg S$ ,  $B \rightarrow S$ , and  $\neg S \wedge (B \vee J)$

# Problem formalization (2)

1. Write a truth table for the three testimonies.

	$B$	$J$	$S$	$J \wedge \neg S$	$B \supset S$	$\neg S \wedge (B \vee J)$
(1)	$T$	$T$	$T$	$F$	$T$	$F$
(2)	$T$	$T$	$F$	$T$	$F$	$T$
(3)	$T$	$F$	$T$	$F$	$T$	$F$
(4)	$T$	$F$	$F$	$F$	$F$	$T$
(5)	$F$	$T$	$T$	$F$	$T$	$F$
(6)	$F$	$T$	$F$	$T$	$T$	$T$
(7)	$F$	$F$	$T$	$F$	$T$	$F$
(8)	$F$	$F$	$F$	$F$	$T$	$F$

# Problem formalization (2)

- Use the truth table to answer the following questions:
  - a) Are the three testimonies satisfiable?  
[Yes, assignment (6) makes them all true]
  - b) The testimony of one of the suspects follows from that of another. Which from which?  
[ $J \wedge \neg S \models \neg S \wedge (B \vee J)$ ]
  - c) Assuming that everybody is innocent, who committed perjury?  
[Everybody is innocent corresponds to assignment (8), and in this case the statements of Brown and Smith are false.]
  - d) Assuming that all testimonies are true, who is innocent and who is guilty?  
[Assuming that all testimonies are true corresponds to assignment (6). In this case Jones is guilty and the others are innocents.]

# Tableaux

Q. Prove the following:  $\models (q \vee p) \rightarrow (p \vee q)$

NOTE: “ $\supset$ ” is a (widely) used alternative representation of “ $\rightarrow$ ”

A. We demonstrate that the negation is not satisfiable (= there is no model satisfying the negated thesis)

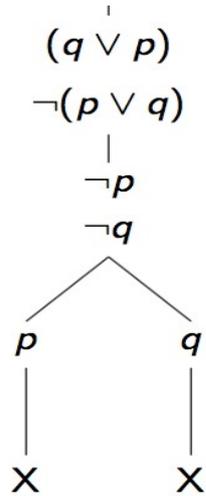
All branches closed.

Rules of the thumb:

- “and”: same node; “or” branches
- transform formulas so that truth is preserved
- In general, prefer “and”s over “or”s

# Tableaux

$\neg(q \vee p \supset p \vee q)$



1. not (q or p) imp (p or q)

1.1 (q or p)  
not (p or q)

1.1.1 not p  
not q

1.1.1.1 p

1.1.1.2 q

1.1.1.1.1 CLOSED

1.1.1.2.1 CLOSED

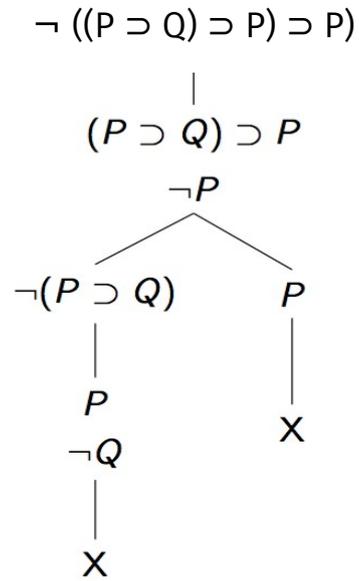
# Tableaux

Q. Prove the following:  $\models (P \rightarrow (Q \rightarrow P)) \rightarrow P$

A. We demonstrate that the negation is not satisfiable (= there is no model satisfying the negated thesis)

All branches closed.

# Tableaux



1. not ((p imp q) imp p) imp p)

1.1 (p imp q) imp p  
not p

1.1.1 not (p imp q)

1.1.2 p

1.1.1.1 p  
not q

1.1.2.1 CLOSED

1.1.1.1.1 CLOSED

# Concluding Remarks

- Prepare on the concepts presented during the course
- Tableaux can be superior to truth tables (consider long formulas or formulas with many different propositional symbols having the same structure)

# Questions:

<https://github.com/avillafiorita/cl-2020>

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